

Zero / Low Cost Counters and Linux TCP for ML

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Zero (or Extremely Low) Cost Counters for Userpace Monitoring



Fast, low overhead, end-to-end metrics to userspace

- Hardware to networking stack and back (ie., Rx and Tx)
- Lockless, no atomics, no memory barriers, no copies
- Always enabled counters
- Constant overhead for a counter irrespective of number of readers
 - e.g., same overhead for 10 readers as with 1 reader
- Inconsistency in a single read is ok with eventual consistency
 - e.g, byte counter and packet counter out of sync on a given read

Flexible Pull of Statistics

- per flow
- per queue
- per device
- in any combination

Stats from networking stack, vendor unique stats in driver, vendor unique stats in h/w

Current APIs for Networking Stats



ethtool -S (or directly invoking ethtool uAPI)

- Vendor / driver unique stats (e.g., per queue, hardware)
- All or nothing scheme (ie., all stats tracked by driver are returned)
- Overhead on every read: rtnl, generating response, copying to userspace, parsing in userspace

ip -s

- Standard device tracked stats (rtnl_link_stats{,64})
- Overhead: same as ethtool

/proc, /sysfs files

• Overhead: system call, data rendering (snprintf), copy to userspace, parsing in userspace

Outline of Infrastructure



Extend netdev-genl or add new stats focused genl

QUERY command to get list of supported counters for specific kernel + driver + H/W

- Returns name and description for each counter group
- Returns name, description, datatype, and optional collection arguments for each counter

REGISTER command for userspace to opt-in to counters using a descriptive name

- TCP socket: tcp/<socket-id>
- L3 stats: ip, ip6
- core stats: core/page_pool
- qdisc stats:
- Driver stats: <netdev>/queue/rx/<id>
- H/W stats: <netdev>/hwqueue/tx/<id>/<collection options>

Names and collection options map to nested attributes

Outline of Infrastructure, cont'd



Counters are mapped to process as read-only page(s)

Successful response returns:

- Base address of read-only mapping
- For each counter: string with a name, data type enum, offset within the mapping

Single counter instance in kernel is observable by many processes - ie., constant cost

- Counters are updated once by kernel
- Counters are read with zero system cost

UNREGISTER command to remove mapping

Notifications to userspace if memory is unmapped by kernel operation

- e.g., monitored socket is closed, netdevice is deleted
- Race here, so may have to convert mapping to zero page to avoid faults

Stat Groups



Existing networking stack statistics moved to a separate page

- All stats in the page are related and mapped as a group
- socket stats, IP, core networking layer, qdisc
- page_pool / generic buffer pool stats

Core code has direct access to stats handlers exported by each layer

Example:

• TCP stats in tcp_sock

struct	tcp stat		
	u64		<pre>/* RFC4898 tcpEStatsAppHCThruOctetsReceived * sum(delta(rcv_nxt)), or how many bytes * were acked. */</pre>
	u32	segs_in;	<pre>/* RFC4898 tcpEStatsPerfSegsIn * total number of segments in. */</pre>
	u32	data_segs_in;	<pre>/* RFC4898 tcpEStatsPerfDataSegsIn * total number of data segments in. */</pre>
	u32 u32 u32 u32 u32	<pre>rcv_nxt; copied_seq; rcv_wup; snd_nxt; segs_out;</pre>	<pre>/* What we want to receive next</pre>
	u32	data_segs_out;	<pre>/* RFC4898 tcpEStatsPerfDataSegsOut * total number of data segments sent. */</pre>
	u64	bytes_sent;	/* RFC4898 tcpEStatsPerfHCDataOctetsOut * total number of data bytes sent. */
	u64	bytes_acked;	<pre>/* RFC4898 tcpEStatsAppHCThruOctetsAcked * sum(delta(snd_una)), or how many bytes * were acked. */</pre>
	u32	dsack_dups;	/* RFC4898 tcpEStatsStackDSACKDups * total number of DSACK blocks received */
;	u32	snd_una;	/* First byte we want an ack for */

S

```
struct tcp_sock {
    ...
    struct tcp_stats *stats; // page allocation that can be mapped
};
```

Driver and Hardware Counters



Stats per queue maintained by driver

- Stats per queue from hardware
- System level stats from hardware

New ndo indicates support for new infrastructure

Hardware stats can have collection options to indicate how counters are accessed

• e.g., updated by timer every N seconds

Support for hardware dependent options

• e.g., hardware push vs software pull to update stats in the page

Depending on H/W details and support - ability to map BAR space to process for some counters

Derived Counters



Computed values

• Callback handler

Timer to update - collection option

Selectable Counters



Ultimate flexibility (with accompanying complexity)

Userspace opts into specific counters

Infrastructure

- Allocates a page of memory for counters
- Assigns slots in the page as counters are added and removed
- Updates stat location to do accounting at a given address
 - 'struct some_object { unsigned int counter; }' -> 'struct some_object { unsigned int *counter; }'
 - Set counter address

More refined but limited to app-specific structs



ebpf allows custom counters across the stack with data going to a userspace map

• Tracepoint + kprobe or various networking hooks

Requires multiple attach points to get end-to-end accounting

Too much overhead for high performance datapath

- Enable tracepoint, probe or bpf attach hook + run bpf program + write to map
- Tracepoints, probes and bpf hook points are code based, not flow based
 - Either limited to 1 system wide stats daemon or take overhead of multiple programs on a tracepoint
- Nanoseconds matter for > 100Gbps flows

Need new hooks to get hardware or driver stats

• See the many, many discussions on ebpf and hardware offloads

Many of the counters duplicate existing system accounting

• e.g., Counters bumped multiple times depending on design



Linux TCP for ML Use Cases

RFC Summary



mempool for huge pages

• expands to custom memory provider

Rx ZC direct to GPU

- devmem patch set packet payload written direct to GPU memory
- based on the memory providers

io_uring support

• extends devmem set to host memory

genl to manage H/W queues

• requirements around use of userspace and GPU memory and flushing references

What do people believe is the end goal of this approach?

• Still significant overhead in the datapath which is not relevant for GPU direct use case

ZC Direct to/from GPU



Should TCP window consider buffers available in buffer pool?

- Data will only land in flow specific buffers
- Slow application == no buffers == dropped packets == waste of network resources

Flushing S/W queues when devmem or host memory is invalidated

- retransmit, ofo and write queues
- references to released (or being released) memory

H/W Queues in Userspace

- Allow application to post buffers directly to hardware
 - Avoids unneeded system calls and page pool overhead proposed in current design
 - Process knows what buffers are available

Other



Big TCP for IPv6

- Drop use of extension header
- Added at L3, removed by driver or needs H/W support
 - Really only needed for tcpdump and it has a solution now
- Make IPv6 equivalent to IPv4 design

Thank You